

$$m \quad a \quad \rightarrow \quad a = m \cdot q + r \quad \text{con } 0 \leq r < m$$

$$a \equiv r \pmod{m} \iff a - r \text{ \u00e9 divisibile per } m$$

$$a \equiv r \pmod{m} \quad b \equiv r' \pmod{m}$$

$$a + b \equiv r + r' \pmod{m} \quad a - b \equiv r - r' \pmod{m}$$

$$a \cdot b \equiv r \cdot r' \pmod{m} \quad \rightarrow \quad / \quad ?? \quad \text{DOPO}$$

$$\text{MOD} = 1000000007;$$

$a \% \text{MOD}$   $\Leftarrow$  resto della divisione per MOD

$$\text{MOD} \% \text{MOD} = 0$$

$$-1 \% \text{MOD} = -1$$

$$\left( (a \% \text{MOD}) + \text{MOD} \right) \% \text{MOD}$$

$$[-\text{MOD} + 1, \text{MOD} - 1]$$

$$[1, 2\text{MOD} - 1]$$

$$\% \text{MOD} \rightarrow 0 \leq x < \text{MOD}$$

$a \cdot b$  NON sta necessariamente in int  $\rightarrow$  long long!

$$e = p_1^{e_1} \dots p_k^{e_k}$$

$$12 = 2^2 \cdot 3$$

$$128 = 2^7$$

- $e$  è un primo
- oppure  $(\min p_i)^2 \leq e$

for (int  $i=2$ ;  $i*i \leq e$ ;  $i++$ ) {

if ( $e \% i == 0$ ) {

//  $i$  è un fattore primo di  $e$

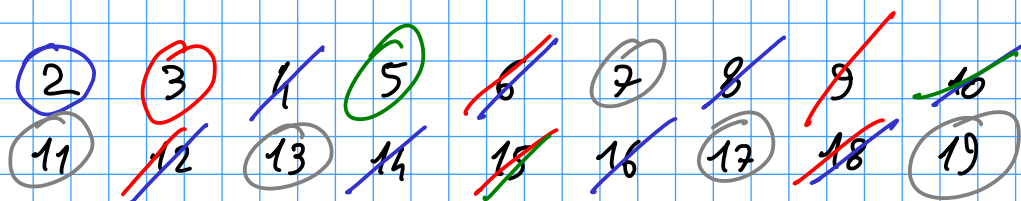
$e /= i$ ;

$i$  -- ;

}

}

//  $e$  è primo o 1



$$\frac{n}{2} + \frac{n}{3} + \frac{n}{5} + \frac{n}{7} \dots$$

$$\rightarrow \sum_{p \leq \sqrt{n}} \frac{n}{p} \leq n \sum_{i \leq \sqrt{n}} \frac{1}{i}$$

$$= O(n \log n)$$

MCD - GCD

MCM - LCM

$\text{GCD}(a, b) = d$  d.c.  $a \% d = 0$   $b \% d = 0$   $d$  è il massimo num. con questa proprietà.

$(6, 8) = 2$      $(11, 7) = 1$      $(2, 4) = 2$

$a = p_1^{a_1} p_2^{a_2} \dots p_n^{a_n}$      $b = p_1^{b_1} p_2^{b_2} \dots p_n^{b_n}$

$(a, b) = p_i^{\min(a_i, b_i)}$

$a = \underline{2}^7 \cdot 5^3 \cdot 11^2$      $b = \underline{2}^6 \cdot 3^4 \cdot 7^{12}$

$d = 2^6 \cdot 3^0 \cdot 5^0 \cdot 7^0 \cdot 11^0 = 2^6 = 64$

$(a, b) = (a, b + k \cdot a) \quad \forall k \in \mathbb{Z}$

$a = n \cdot d$      $b + k \cdot a = (n + k \cdot m) \cdot d$   
 $b = m \cdot d$

$(a, b) = (a, b + k \cdot a - k \cdot a)$

$(a, 0) = a$

```
int gcd(int a, int b) {
  if (a < b) return gcd(b, a);
  if (b == 0) return a;
  return gcd(b, a % b);
}
```

$a = kb + r \rightarrow r = a - kb$

$O(\log a)$

$(a, b) = 1$  Coprimi

Dato  $x$ , esiste  $a$  :  $ax = 1 \pmod{m}$  ?

$$\left( \begin{array}{l} b/x \pmod{m} \text{ def: } (b/x) \cdot x \equiv b \pmod{m} \\ \parallel \\ b \cdot a \qquad \qquad b \cdot a \cdot x \equiv b \cdot (a \cdot x) \equiv b \pmod{m} \end{array} \right)$$

$\Leftrightarrow (m, x) = 1$  Se  $m$  è primo  $p \Leftrightarrow x \not\equiv 0 \pmod{p}$

Teorema di Fermat

$$x^{p-1} = 1 \pmod{p} \text{ se } x \not\equiv 0 \pmod{p}$$

$$x^{p-2} \cdot x = x^{p-1} = 1 \pmod{p}$$

$$a = x^{p-2} \Rightarrow ax = 1$$

$$O(\log p)$$

$\phi(m)$  = # di numeri  $\leq m$  coprimi con  $m$

$$\phi(p) = p-1 \quad \phi(p^k) = p^{k-1}(p-1)$$

$$\phi(a \cdot b) = \phi(a) \cdot \phi(b) \text{ se } (a, b) = 1$$

Teorema di Eulero

$$x^{\phi(m)} = 1 \pmod{m} \text{ se } (x, m) = 1$$

$$x^{\phi(m)-1} \rightarrow \text{inverso di } x \pmod{m}$$

$\binom{n}{k} = \#$  di sottoinsiemi di dimensione  $k$  in un insieme con  $n$  elem.

$$\frac{n!}{k!(n-k)!}$$

Contiamo  $\#$  di sottoinsiemi ordinati di  $k$  elementi in  $n$

•  $k \cdot (k-1) \cdot (k-2) \dots 1 \cdot \binom{n}{k} = k! \binom{n}{k}$

•  $n \cdot (n-1) \dots (n-k+1) = \frac{n \cdot (n-1) \dots (n-k+1) \cdot (n-k) \dots 1}{(n-k) \dots 1}$

$$= \frac{n!}{(n-k)!}$$

$$k! \binom{n}{k} = \frac{n!}{(n-k)!} \Rightarrow \binom{n}{k} = \frac{n!}{k!(n-k)!}$$

$$\binom{n}{k} = \frac{n \cdot (n-1) \cdot (n-2) \dots (n-k+1)}{k!}$$

- calcolare  $\binom{n}{0} \binom{n}{1} \dots \binom{n}{n}$

$$\binom{n}{k+1} = \frac{n!}{(k+1)!(n-k-1)!} = \frac{n!}{k!(n-k)!} \cdot \frac{n-k}{k+1} = \frac{n-k}{k+1} \binom{n}{k}$$

- calcolare  $\binom{k}{k} \binom{k+1}{k} \dots \binom{n}{k}$

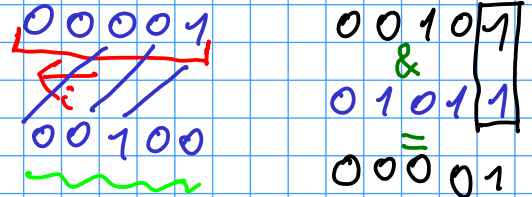
$$\binom{n}{k} = \frac{n!}{k!(n-k)!} = \frac{n}{n-k} \cdot \frac{(n-1)!}{k!(n-k-1)!} = \frac{n}{n-k} \binom{n-1}{k}$$

-  $\binom{n}{0} + \dots + \binom{n}{n} = 2^n$

$\{a, b, c, d, e\}$

$01001 \Leftrightarrow \{b, e\}$

(a) controllare appartenenza:  $i \in A \Leftrightarrow (\underline{1ULL \ll i}) \& A \neq 0$



- se l' $i$ -esimo bit di  $A$  è 1,   
 - altrimenti, 0

(b) intersezione  $A, B$  :  $C = A \cap B \Leftrightarrow C = A \& B$

(c) unione  $A, B$  :  $C = A \cup B \Leftrightarrow C = A | B$

01001

11010

11011

(d)  $A \leftarrow A \cup \{i\}$       $A | (\underline{1ULL \ll i})$

(e)  $A^c$       $\sim A$       $\sim 01001 = 10110$

(f)  $A \cap B$       $A \& (\sim B)$

(g) togliere un elemento  $A \& \sim(1ULL \ll i)$

(h) differenza simmetrica  $A \Delta B$       $A \wedge B$

01001

10011

11010

(i) cambiare (toggle) un elemento

$A \wedge (1ULL \ll i)$

$$0x1F00 \Rightarrow \underline{00011110} \underline{00000000}$$

$$0b01011100$$

$$A \& (\sim(\sim 0ULL) \ll 5)$$

$$111...100000$$

$$\underline{0000...011111}$$

$$|A| = \text{--BUILTIN\_POPCOUNTLL}(A);$$

$$n \& -n \quad -n = (\sim n) + 1$$

ESERCIZIO

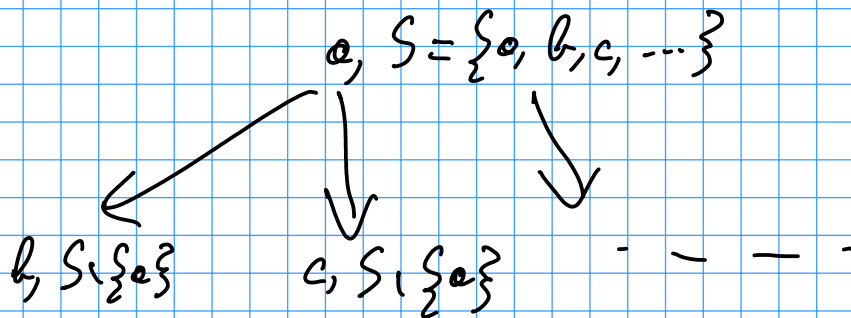
COMMESSE VIAGGIATORE

$G$ ,  $n$  nodi, completo, pesati

trovare il ciclo con  $n$  nodi di costo minimo

Esiste algo.  $O(n^{\binom{n}{2}}) \ll O(n!)$

Q: per ogni sottoinsieme  $S$  di nodi, quale è il cammino di costo minimo che li visita tutti, parte da  $a \in S$  e arriva in  $a$ ?



$$a, \{a\} \rightarrow w(a, a)$$

$$\sum_{A \subseteq \{1, \dots, n\}} |A|^2 = n^2 2^n$$

- for (v. long long sub = 0; sub < (1ULL << n); sub++) ...

$$1ULL \& (A == 0)$$

{ 1, ..., n }    1 2 3 4     $\rightarrow$     4 2 1 3

STB:; NEXT - PERMUTATION

VECTOR<INT> PERM(n);

iota(PERM.begin(), PERM.end(), 0);

do {

...

} while (next\_permutation(PERM.begin(), PERM.end()));

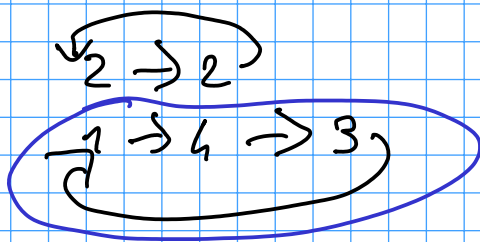
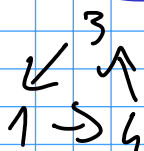
PERM, INVERSA;

•  $INV\_PERM[PERM[i]] = i;$

$PERM[INV\_PERM[i]] = i;$

DECOMP. CICLI

1 2 3 4  $\rightarrow$  4 2 1 3



2)

$PERM[] = \{4, 2, 1, 3\};$

$1 \rightarrow 4$